



Optimizing the RJMS of an Academic HPC & Research Computing Facility

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Summary



2 Toward Improved User Job Management through a Novel RJMS configuration

③ Impact Analysis of the Updated Models and Policies through Workload Performance Evaluation





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Introduction: Context and Motivations

Summary

1 Introduction: Context and Motivations

2 Toward Improved User Job Management through a Novel RJMS configuration

3 Impact Analysis of the Updated Models and Policies through Workload Performance Evaluation

4 Conclusion & Perspectives



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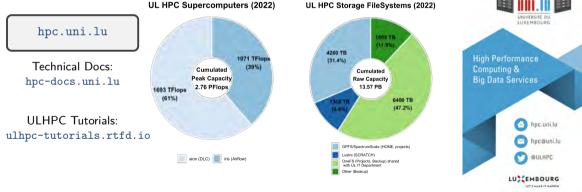


Introduction: Context and Motivations

Uni.lu HPC (UL HPC) Facility

Managed and operated since 2007 (Dr. S. Varrette & Co.)

 → 2nd Largest HPC facility in Luxembourg after EuroHPC MeluXina

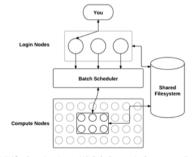




Resource and Job Management Systems

• Resource and Job Management System (RJMS)

- \hookrightarrow Glue for a parallel computer to execute parallel jobs
- $\,\hookrightarrow\,$ Goal: satisfy users demands for computation
 - $\checkmark\,$ assign resources to user jobs with an efficient manner



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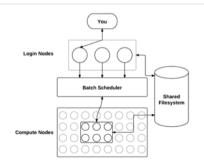
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• HPC Resources:

- $\,\hookrightarrow\,$ Nodes (typically a unique IP address)
 - ✓ Sockets / Cores / Hyperthreads
 - ✓ Memory
 - $\checkmark~$ Interconnect/switch resources
- \hookrightarrow Generic resources (e.g. GPUs)
- $\, \hookrightarrow \, \, {\sf Licenses}$
- Strategic Position
 - $\, \hookrightarrow \, \, \mathsf{Direct}/\mathsf{constant} \, \, \mathsf{knowledge} \, \, \mathsf{of} \, \mathsf{resources}$
 - $\,\hookrightarrow\,$ Launch and otherwise manage jobs



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https://slurm.schedmd.com/

- ULHPC uses Slurm for cluster/resource management and job scheduling
 - \hookrightarrow Simple Linux Utility for Resource Management
 - $\,\hookrightarrow\,$ Reference RJMS serving most of Top500 systems
 - ✓ official documentation, tutorial, FAQ







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Slurm on ULHPC clusters

- ✓ official documentation, tutorial, FAQ
- Seminal configuration part of flagship iris production release (in 2017)
 - $\,\hookrightarrow\,$ migration from OAR RJMS
 - $\hookrightarrow \mbox{Slurm provides superior scalability and performance [JSSPP12], inherent compatibility with multiple distributed libraries (Dask, IPyparallel...) and MPI suits (though PMI[x])$
 - \checkmark brings a more convenient and flexible interface for non-specialists
 - \hookrightarrow inspired from other HPC centers (LLNL, Niflheim, CSCS...) & Simulators [PEARC18]

[JSSPP12] Y. Georgiou, M. Hautreux "Evaluating scalability and efficiency of the resource and job management system on large HPC clusters", in W. on Job Scheduling Strategies for Parallel Processing (LNCS JSSPP'12), Springer, pp. 134–156 (2012). [PEARC18] N. A. Simakov & al. "Slurm Simulator: Improving Slurm Scheduler Performance on Large HPC Systems by Utilization of Multiple Controllers and Node Sharing" in Proc. of the ACM Practice and Experience on Advanced Research Computing (PEARC'18) (2018).





After 3 years of production on iris...

(1/2)

- Over-complexified setup for partition/QOS
 - \hookrightarrow integration of hetegeneous hardware (GPU nodes, Large-memory) starting 2018





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- Over-complexified setup for partition/QOS
 - \hookrightarrow integration of hetegeneous hardware (GPU nodes, Large-memory) starting 2018
 - \hookrightarrow under vs. over used partitions
 - \checkmark batch partition (Dual-CPU regular nodes) saturated and over-used
 - \checkmark dedicated resources for interactive or long (regular nodes only) under-used



(1/2)



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 - \hookrightarrow under vs. over used partitions
 - \checkmark batch partition (Dual-CPU regular nodes) saturated and over-used
 - \checkmark dedicated resources for interactive or long (regular nodes only) under-used
 - \hookrightarrow for each partition, associated QOS queue named qos-<partitionname>[-XXX]
 - $\checkmark\,$ allows for specific research groups/industrial partners privileged/exclusive access
 - ✓ Ex: qos-batch (default), qos-batch-[001-...] (Group Prof. / Partner $XX_{\{1-..\}}$)
 - ✓ Ex: qos-gpu (default), qos-gpu-[001-...] (Group Prof. / Partner YY {1-...})
 - \checkmark Ex: qos-covid ultra high priority jobs to support fight against COVID-19
 - $\,\hookrightarrow\,$ In practice, this setup brought frustration, jealousy & confusion within users



(1/2)



After 3 years of production on iris...

(2/2)

- No cross-partition QOS except best-effort (preemptible jobs):
 - \hookrightarrow interactive jobs on non-regular nodes (GPU, large-memory) artificially complexified
 - \hookrightarrow each specific usage treated by dedicated QOS (Ex: qos-batch-0*, qos-covid)





After 3 years of production on iris...

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- No cross-partition QOS except best-effort (preemptible jobs):
 - \hookrightarrow interactive jobs on non-regular nodes (GPU, large-memory) artificially complexified
 - \hookrightarrow each specific usage treated by dedicated QOS (Ex: qos-batch-0*, qos-covid)
 - ✓ no global priority level (low→urgent), unreadability of QOS objectives vs. partition
- Fairsharing relying on the Depth-Oblivious Fair-share Factor algorithm
 - \hookrightarrow variant of classical fair-share factor, increases usable priority ranges
 - \checkmark very complex algorithm, hard to explain impacts/issues with job priority
 - $\,\hookrightarrow\,$ no policy for raw-share attribution lead to unfair situations
 - \hookrightarrow hard to evaluate/size associated Slurm parameters Ex: PriorityWeight* (among others)





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• Incomplete Account hierarchy

- \hookrightarrow not able to cover new workload requests associated to funded projects/training events
- \hookrightarrow novel auditing capabilities on the platform usage / cost model requested
 - \checkmark \ldots and need to be integrated in the Slurm configuration



(2/2)



Toward a Novel RJMS Setup

- Acquisition & integration of new liquid-cooled supercomputer aion
 - $\,\hookrightarrow\,$ occasion to deeply review and ${\bf optimize}$ the seminal configuration
 - $\,\hookrightarrow\,$ mitigate the identified pitfalls & take advantage of experience gained







Toward a Novel RJMS Setup

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In this talk: ULHPC Slurm configuration 2.0

- 1 Re-defining the partition, QOS and accounting model \hookrightarrow offer a more focused and easy-to-use configuration
- Review fair-sharing model giving **incentives** to good practices
 - $\,\hookrightarrow\,$ take into account monetary contribution to increase priority
 - \hookrightarrow tribute to past efficient usage
- Rethinking global resources limits set to the partitions / accounts associations / QOS
- Consolidating the RJMS setup for HA services
- **5** Define common [federated vs. multi-cluster] scheduling / accounting policy











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ULHPC Slurm Partitions 2.0 -p, -partition=<partition>

- Tied to global types/classes of available computing nodes batch[,gpu][,bigmem]
- New floating partition interactive across all nodes for (short and quick) tests
 - $\hookrightarrow \text{ selection of the expected resource type through feature} \qquad \texttt{-C broadwell,skylake,gpu,volta[32]} \ldots$
 - $\, \hookrightarrow \,$ backfill scheduling enabled and optimized to favor interactive and/or small jobs if queued
- Max wall clock time for user job on non-floating partition reduced from 5 to 2 days

les	MaxNodes	MaxTime	DefaultTime	PriorityTier	#Node	Туре	AION Partition
TED	UNLIMITED	5d	10h	100	318	hidden	all (sysadmins)
	2	2h	30min	100	318	floating	interactive
	64	48h	2h	1	318	-	batch
les	MaxNodes	MaxTime	DefaultTime	PriorityTier	#Node	Туре	IRIS Partition
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	2	2h	30min	100	196	floating	interactive
	64	48h	2h	1	168		batch
		48h	2h	1	24		gpu
	4	4011	211	1	24		

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ULHPC Slurm QOS 2.0

- $\bullet~\mbox{New Cross-partition QOS},$ mainly tied to priority level (low $\rightarrow~\mbox{urgent})$
 - $\,\hookrightarrow\,$ Simpler names than before (i.e. no more qos- prefix)
 - $\,\hookrightarrow\,$ special preemptible QOS kept for best-effort jobs: besteffort
 - $\,\hookrightarrow\,$ new long QOS allows to run jobs for up to 14 days (instead of the default 2 days)
- Further limits on Slurm Trackable RESources (TRES)

QOS	Partition	Allowed [L1] Account	Prio	GrpTRES	MaxTresP.	J MaxJobPU	Flags
besteffort	*	ALL	1			100	NoReserve
low	*	ALL (default for CRP/externals)	10			2	DenyOnLimit
normal	*	Default (UL, Projects,)	100			50	DenyOnLimit
long	*	UL,Projects,etc.	100	node=12	node=2	4	DenyOnLimit, PartitionTimeLimi
debug	interactive	ALL	150	node=8		2	DenyOnLimit
high	*	(restricted) UL, Projects, Industry	200			10	DenyOnLimit
urgent	*	(restricted) UL, Projects, Industry	1000			100	DenyOnLimit
admin	all	(restricted) sysadmins	1000				DenyOnLimit







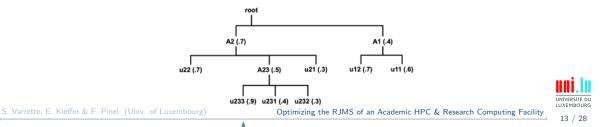
ULHPC Fairsharing 2.0

- Three fairsharing algorithms implemented in Slurm:
 - $\,\hookrightarrow\,$ Classic Fairshare, Depth-Oblivious Fair-share (initial setup) and Fair Tree
 - $\,\hookrightarrow\,$ Thourough evaluation of all 3 fairshare algorithms (Python-based simulator)



ULHPC Fairsharing 2.0

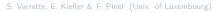
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 - $\,\hookrightarrow\,$ Thourough evaluation of all 3 fairshare algorithms (Python-based simulator)
- In Fair Tree, all users from higher priority account receive higher fair share factor
 - $\,\hookrightarrow\,\ldots\,$ when compared to all users from a lower priority account
 - $\,\hookrightarrow\,$ Done with rooted plane tree (rooted ordered tree)
 - $\checkmark\,$ logically created then sorted by fairshare level with highest fairshare values on the left
 - \checkmark tree is then visited in a depth-first traversal way





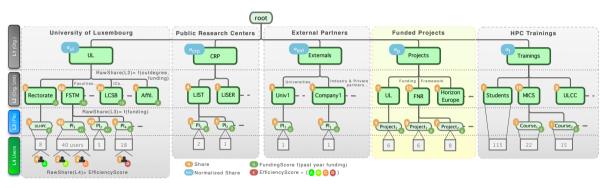
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 - $\,\hookrightarrow\,$ Classic Fairshare, Depth-Oblivious Fair-share (initial setup) and Fair Tree
 - $\,\hookrightarrow\,$ Thourough evaluation of all 3 fairshare algorithms (Python-based simulator)
- New configuration with Multifactor Priority Plugin and Fair tree algorithm
 - \hookrightarrow efficiency: new jobs are immediately assigned a priority
 - \hookrightarrow fairshare levels are more easily understandable
 - $\,\hookrightarrow\,$ YET quite sensitive to the raw shares associated to each user account
- Necessity to deeply restructure associations & shares in the accounting DB
 - $\,\hookrightarrow\,$ formalize consistent rules to attribute raw shares
 - $\,\hookrightarrow\,$ give novel incentives to good practices
 - $\checkmark~$ take into account monetary contribution to increase priority
 - $\checkmark~$ tribute to past efficient usage





Account Hierarchy 2.0





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Efficiency Score (L4)

- Updated every year based on past jobs efficiency.
 - \hookrightarrow Similar notion of "nutri-score': A(very good 3), B (good: 2), C (bad, 1), D(very bad 0)
- Proposed Metric for user U: Average Wall-time Accuracy (WRA) (higher the better)
 - \hookrightarrow Defined for a given time period (past year)

 $\texttt{sacct} - \texttt{u} < \texttt{U} > -\texttt{X} - \texttt{S} < \texttt{start} > -\texttt{E} < \texttt{end} > \texttt{[...]} \ \# \ --format \ \textit{User, JobID, state, time, elapsed}$

 $egin{aligned} S_{ ext{efficiency}}(U, Year) &= \mathsf{WRA}(U, Year) \ &= rac{1}{N} \sum_{JobID} rac{T_{ ext{elapsed}}(JobID)}{T_{ ext{asked}}(JobID)} \end{aligned}$

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• U raw share: $1 + S_{\text{efficiency}}(U, Year)$

Score Avg. W			
A (3) very good B (2) good C (1) bad D (0) very bad	$\begin{array}{r} S_{\rm efficiency} \geq 75\% \\ 50\% \leq S_{\rm efficiency} < 75\% \\ 25\% \leq S_{\rm efficiency} < 50\% \\ S_{\rm efficiency} < 25\% \end{array}$		



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• WIP: integrate other efficiency metrics (CPU, mem, GPU efficiency)



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Job Accounting and Billing

- Utilization of the University computational resources is charged in Service Unit (SU)
 - $\,\hookrightarrow\,$ 1 SU $\simeq 1$ hour on 1 physical processor core on regular computing node
 - \hookrightarrow Usage charged 0,03€ per SU (VAT excluded) (external partners, funded projects etc.)
- A job is characterized* (and thus billed) according to the following elements:
 - \hookrightarrow T_{exec} : Execution time (in hours)
 - $\,\hookrightarrow\,$ $\textit{N}_{Nodes}:$ number of computing nodes, and per node:
 - \checkmark N_{cores}: number of CPU cores allocated per node
 - $\checkmark~$ Mem: memory size allocated per node, in GB
 - \checkmark N_{gpus}: number of GPU allocated per node
 - \hookrightarrow associated weighted factors $\alpha_{\textit{cpu}}, \alpha_{\textit{mem}}, \alpha_{\textit{GPU}}$ defined as TRESBillingWeight in Slurm
 - $\checkmark\,$ account for consumed resources other than just CPUs, taken into account in fairshare factor
 - \checkmark α_{cpu} : normalized relative perf. of CPU processor core (reference: skylake 73,6 GFlops/core)
 - $\checkmark ~ \alpha_{\it mem}$: inverse of the average available memory size per core
 - $\checkmark ~ \alpha_{\textit{GPU}}$: weight per GPU accelerator





Job Accounting and Billing

Number of SU associated to a job

 $N_{\text{Nodes}} \times [\alpha_{cpu} \times N_{\text{cores}} + \alpha_{mem} \times Mem + \alpha_{gpu} \times N_{gpus}] \times T_{\text{exec}}$

• Current billing weights:

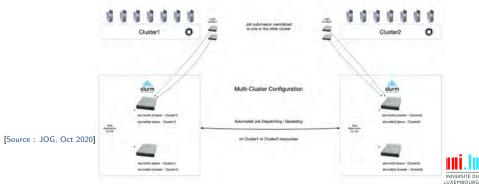
Cluster	Node Type	Partition	#Cores/node	CPU	$lpha_{ m cpu}$	α_{mem}	lpha GPU
Iris, Aion	Regular	interactive	28/128	n/a	0	0	0
Iris	Regular	batch	28	broadwell	0.522	$\frac{1}{4} = 0,25$	0
Iris	Regular	batch	28	skylake	1.0	$\frac{1}{4} = 0,25$	0
Iris	GPU	gpu	28	skylake	1.0	$\frac{1}{27}$	50
lris	Large-Mem	bigmem	112	skylake	1.0	$\frac{1}{27}$	0
Aion	Regular	batch	128	ерус	0,57	$\frac{\frac{2}{1}}{1.75}$	0





Federated vs. Multi-Cluster Scheduling Strategy

• Two possible approaches to integrating aion into the existing RJMS configuration \hookrightarrow Multi-Cluster: allow to submit jobs across each clusters sbatch -M {iris,aion} [...]



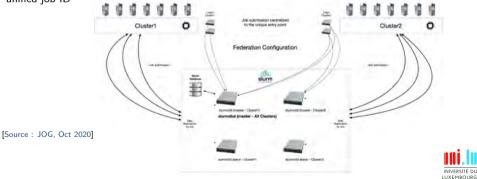
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 - $\,\hookrightarrow\,$ Federation: P2P replication of job submitted locally, distributed schedule attempt
 - ✓ unified job ID





Federated vs. Multi-Cluster Scheduling Strategy

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 - $\hookrightarrow\,$ Federation: P2P replication of job submitted locally, distributed schedule attempt $\checkmark\,$ unified job ID

Selected Setup: Hybrid Multi-Cluster Scheduling Strategy

- Redundant master/slave RJMS controllers associated to each cluster
- Shared (like for federations), the same Slurm database service (slurmdb)

 → facilitate and centralize accounting information and management



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Summary



 $\fbox{2}$ Toward Improved User Job Management through a Novel RJMS configuration

③ Impact Analysis of the Updated Models and Policies through Workload Performance Evaluation

Impact Analysis of the Updated Models and Policies through Workload Performance Evaluation





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Performance Evaluation and Experimental Setup

• Impact of the updated RJMS policy and setup particularly hard to qualify

- ↔ Approach 1: Slurm simulators of workloads execution. Ex: BSC slurm_simulator [PMBS18]
 - ✓ **YET** old/obsolete version of Slurm featured (17.x)
 - \checkmark novel proposed contributions absent and thus to implement. . .

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Our chance (Approach 3)

- All the modifications/proposals applied at once within our production systems! → on Oct 22, 2020 during a maintenance session.
 - \hookrightarrow 1 supercomputer (iris) has thus seen **both** configs for a sufficient amount of time
 - $\checkmark\,$ workload analysis able to reasonably capture impact of the changes on RJMS perf.

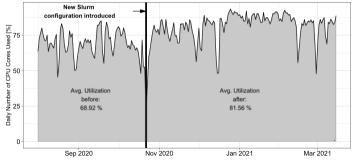




Impact Analysis of the Updated Models and Policies through Workload Performance Evaluation

Impact on the Total Utilization

- Impact of the updated Slurm configuration on the ULHPC relative utilization:
 - $\, \hookrightarrow \,$ aggregates traces from several months of uninterrupted HPC services
 - $\,\hookrightarrow\,$ daily number of CPU cores used increased by 12.64%: \simeq 81.56% over 6 months



UL HPC Facility Relative Utilization (CPUHours)

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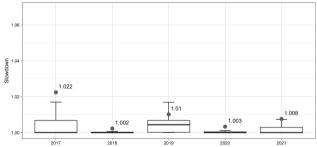
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Impact Analysis of the Updated Models and Policies through Workload Performance Evaluation

Impact on the Job Slowdown

- Slowdown: traditional performance metrics for job scheduling strategies [Feitelson15] slowdown = $\frac{T_w + T_r}{T}$
 - RJMS response time normalized by the runtime: \hookrightarrow
 - \leftrightarrow average yearly slowdown very close to the optimal value (1)



Slowdown evolution on the iris supercomputer

[Feitelson15] D. G. Feitelson, "Workload Modeling for Computer Systems Performance Evaluation, 1st ed. USA: Cambridge University Press (2015

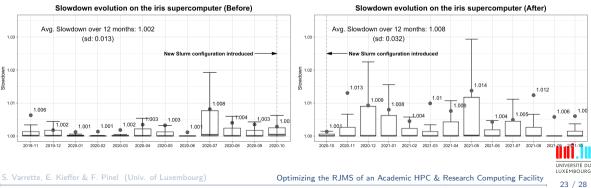
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Impact on the Job Slowdown

- More detailed analysis to better capture the impact of the updated RJMS config.
 - \hookrightarrow Monthly slowdown for all iris jobs run for 1 year period before and after the changes
 - $\,\hookrightarrow\,$ Fairly negligible impact: observed avg. job slowdown only increased by 0.59%

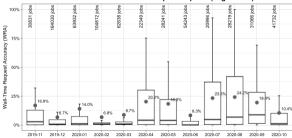




Impact on the Average Wall-time Request Accuracy

• Evaluation covering 1 year period before

proposed configuration change



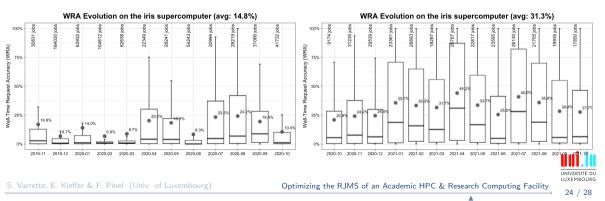
WRA Evolution on the iris supercomputer (avg: 14.8%)





Impact on the Average Wall-time Request Accuracy

- Evaluation covering 1 year period before and after proposed configuration change
- Avg. WRA for the processed jobs was increased by 110.81%
 - $\,\hookrightarrow\,$ moving from 14.8% on average to 31.3%

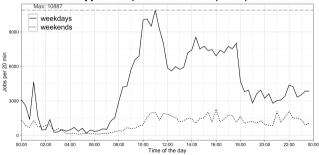




Impact Analysis of the Updated Models and Policies through Workload Performance Evaluation

Current Daily Arrival Pattern

- Proposed changes in production for 18 months and extensively used
 - $\hookrightarrow \text{ Ex: Daily Arrival Pattern extracted from iris workload trace over 2021} \quad \texttt{Jan 1^{st}} \rightarrow \texttt{Dec 31^{th}}$
 - $\checkmark\,$ up to 12869 jobs submitted (incl. 10887 outside weekends) and processed per 20 min



Daily job arrival patterns on the iris supercomputer





Summary



0 Toward Improved User Job Management through a Novel RJMS configuration

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4 Conclusion & Perspectives



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Conclusion

- In this talk:
 - $\,\hookrightarrow\,$ Optimization of Slurm RJMS config when acquiring & integrating a new supercomputer
 - \checkmark smooth integration within the existing HPC ecosystem
 - $\checkmark~$ mitigation of the identified pittfalls of the initial configuration after 3y of production





Conclusion

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$\,\hookrightarrow\,$ Novel and flexible setup, adaptable to other HPC centers

- $\checkmark~$ new scheme defining partitions, QOS queues, priorities & their resource limits
- \checkmark fairsharing mechanism updated, with **incentives** to good practices & monetary contribution
- \checkmark rededigned account hierarchy, hybrid multi-cluster scheduling strategy





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- $\checkmark~$ rededigned account hierarchy, hybrid multi-cluster scheduling strategy
- $\,\hookrightarrow\,$ New cost model and updated policies in production for 18 months
- $\,\hookrightarrow\,$ Performance evaluation from real workload traces
- Perspectives and Future directions
 - \hookrightarrow smooth integration with Euro-HPC infrastructures
 - \checkmark transparently outsource Research Computing/data analytic workflows to Tier-0 systems
 - $\,\hookrightarrow\,$ model & automatically offload from RJMS some of the less-demanding jobs
 - ✓ target dynamically allocated virtual cloud resources (burst instances)







Questions?

Sebastien Varrette, Emmanuel Kieffer and Frederic Pinel Optimizing the Resource and Job Management System of an Academic HPC & Research Computing Facility – IEEE ISPDC 2022 University of Luxembourg, Belval Campus Maison du Nombre, 4th floor 2, avenue de l'Université L-4365 Esch-sur-Alzette mail: firsname.lastname@uni.lu

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Conclusion & Perspectives

High Performance Computing @ Uni.lu





ULHPC Technical Docs

hpc-docs.uni.lu



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Appendix / Backup Slides



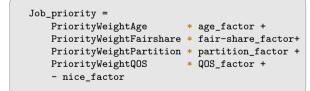
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ULHPC Job Prioritization Factors

- Age: length of time a job has been waiting (PD state) in the queue
- Fairshare: difference between the portion of the computing resource that has been promised and the amount of resources that has been consumed
- Partition: factor associated with each node partition
 - $\,\hookrightarrow\,$ Ex: privilege interactive over batch
- QOS A factor associated with each Quality Of Service (low \longrightarrow urgent)







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Job_priority =		
PriorityWeightAge	*	age_factor +
PriorityWeightFairshare	*	fair-share_factor+
PriorityWeightPartition	*	<pre>partition_factor +</pre>
PriorityWeightQOS	*	QOS_factor +
- nice_factor		
	PriorityWeightAge PriorityWeightFairshare PriorityWeightPartition PriorityWeightQOS	PriorityWeightAge * PriorityWeightFairshare * PriorityWeightPartition * PriorityWeightQOS *

# Current weights on ULHPC platform						
\$ sprio	-w # -	-format "%8	8i %5A %9F	%9P %Q"		
JOBID	AGE	FAIRSHARE	PARTITION	QOS		
Weights	2000	3000	10000	1000		





Account Hierarchy 2.0

Accounting records re-organized as a hierarchical tree (3 layers L_{1,2,3} + leafs)
 → L1: Organization Level: UL, CRPs, Externals, Projects, Trainings

 ✓ guarantee 85% of the shares for core UL activities
 → L2: Organizational Unit (Faculty, ICs, External partner, Funding program...)
 ✓ raw share depends on outdegree and funding score
 → L3: Principal Investigator (PIs), Projects, Course
 ✓ raw share depends on funding score (different weight)
 ✓ eventually restricted only to projects and courses
 → L4: End User (ULHPC login)
 ✓ Raw share based on efficiency score





Funding Score (L2/L3)

- Associated with account A belonging to level L in the hierarchy
 - $\,\hookrightarrow\,$ yearly updated at the beginning of the year
 - $\,\hookrightarrow\,$ depreciation based on contribution type, weighted by level threshold $\beta_{\rm L}$

$$ext{FundingScore}_L(A) = \left\lfloor eta_L rac{\textit{Investment}_A(\textit{Year}-1)}{\#\textit{months}}
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floor$$





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- Ex1: Exceptional contribution of $120K \in$ performed in 2021 by a faculty (L2 account A)
 - \hookrightarrow depreciation: 12 months (*default*)
 - \hookrightarrow funding score in 2022: $\left[\beta_{L_2} \frac{120000}{12}\right] = \left[\beta_{L_2} \times 10000\right].$
- Ex2: let P be a project granted in 2021 to start in 2022 for a duration of 36 months
 - \hookrightarrow <code>budget: 27K</code> <code>allocated</code> for HPC costs
 - \hookrightarrow funding score for the years 2022, 2023 and 2024: $\left|\beta_{L_3} \frac{27000}{36}\right| = \left|\beta_{L_3} \times 750\right|$

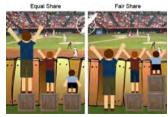


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ULHPC Fairsharing 2.0

- Fairsharing: way of ensuring that users get their appropriate portion of a system
 - $\,\hookrightarrow\,$ Share: portion of the system users have been granted.
 - $\,\hookrightarrow\,$ Usage: amount of the system users have actually used.
 - $\,\hookrightarrow\,$ Fairshare score: value the system calculates based off of user's usage.
 - $\checkmark\,$ difference between the portion of the computing resource that has been promised and the amount of resources that has been consumed
 - $\,\hookrightarrow\,$ Priority score: priority assigned based off of the user's fairshare score.



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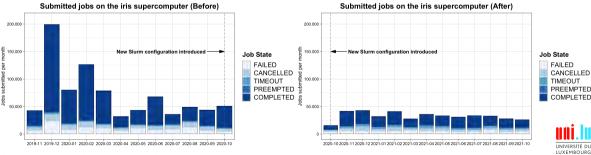




Impact on the Daily submitted Jobs

- Obviously depends on the usage pattern of the platform
 - $\,\hookrightarrow\,$ Avg. #monthly submitted jobs: 71102 \rightarrow 35384
 - $\,\hookrightarrow\,$ Submission pattern changed and hard to conclude
 - ✓ COMPLETED job decreased...
 - ✓ Yet roughly eq. to CANCELLED+TIMEOUT increase (user-dependant)

Job State	Quantity Before	Quantity After	Difference
COMPLETED	73.3%	60.3%	-13%
PREEMPTED	0.12%	0.02%	-0.1%
FAILED	13.1%	15.2%	+2.1%
CANCELLED	10.4%	14.5%	+4.1%
TIMEOUT	3.06%	9.92%	+6,86%



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